

CSCI 6214— Randomized Algorithms— FALL 2021
Problem Set #4

Problem 1:

Consider throwing m balls into n bins, and for convenience let the bins be numbered from 0 to $n - 1$. We say there is a k - gap starting at bin i if bins $i, i + 1, \dots, i + k - 1$ are all empty.

- Determine the expected number of k -gaps.
- Prove a Chernoff-like bound for the number of k -gaps. (Hint: If you let $X_i = 1$ when there is a k - gap starting at bin i , then there are dependencies between X_i and X_{i+1} ; to avoid these dependencies, you might consider X_i and X_{i+k} .)

Problem 2:

Suppose that we vary the balls-and-bins process as follows. For convenience let the bins be numbered from 0 to $n - 1$. There are $\log_2 n$ players. Each player randomly chooses a starting location l uniformly from $[0, n - 1]$ and then places one ball in each of the bins numbered $l \bmod n, l + 1 \bmod n, \dots, l + n/\log_2 n - 1 \bmod n$. Argue that the maximum load in this case is only $O(\log \log n / \log \log \log n)$ with probability that approaches 1 as $n \rightarrow \infty$.

Problem 3:

We consider another way to obtain Chernoff-like bounds in the setting of balls and bins without using the relationship between the real distribution and the Poisson distribution we saw in class. Consider n balls thrown randomly into n bins. Let $X_i = 1$ if the i -th bin is empty and 0 otherwise. Let $X = \sum_{i=1}^n X_i$. Let $Y_i, i = 1, \dots, n$, be independent Bernoulli random variables that are 1 with probability $p = (1 - 1/n)^n$. Let $Y = \sum_{i=1}^n Y_i$.

- Show that $E[X_1 X_2 \dots X_k] \leq E[Y_1 Y_2 \dots Y_k]$ for any $k \geq 1$.
- Show that $E[e^{tX}] \leq E[e^{tY}]$ for all $t \geq 0$. (Hint: Use the expansion for e^x and compare $E[X^k]$ to $E[Y^k]$.)
- Derive a Chernoff bound for $Pr(X \geq (1 + \delta)E[X])$.